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(Extended Abstract)

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Inheritance of Interface Specifications
(Extended Abstract)

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Abstract
Four alternatives for the semantics of inheritance of specifications are
discussed. The information loss and frame axiom problems for inherited
specifications are also considered.

1 Introduction
An interface specification language (ISL) defines both how to call a module and
its (functional) behavior [Win83] [Win87] [Lam89] [GHG+93]. The details of
how to call a module and some aspects its behavior are specific to the particular
programming language; hence in the Larch approach to interface specification
[GHG+93], each ISL is tailored to a particular programming language.
What does this tailoring involve?

• The syntax for specifying interfaces is a subset of the syntax for the pro-
gramming language, so it can be directly compared to the interface of a
candidate implementation. This means that the ISL must use the type
system of the programming language.

• Most of the semantic concepts of the programming language should be
reflected in the ISL’s semantics.

• The ISL should allow the specifier to use the programming language’s ab-
straction mechanisms. For example, an ISL tailored to an object-oriented
programming language (OOPL) should allow inheritance of specifications.

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• The ISL should ease ways of reasoning about the correctness of code that are common or otherwise important. For example, an ISL for an OOPL should ease reasoning that uses supertype abstraction (thinking only about the types written in the program, not about the dynamically possible subtype that expressions may denote [IW90]).

It follows that when one is designing an ISL for an OOPL, one must consider both subtyping and inheritance.

1.1 Modularity

The last point in the list of ways to tailor an ISL is especially important in an object-oriented context. In order to support supertype abstraction one must pay attention to modularity. An ISL is modular if:

• at any point where an argument is specified to have type $T$, an actual argument of some subtype of $T$ is allowed,

• when adding new subtypes of existing types, or when using specification inheritance, one need not change any existing specifications, and the meaning of the existing specifications does not change (except for allowing the new subtypes to be used),

• to specify a type, one need not be aware of the specification of any other types, except those of the supertypes of the type being specified (and their supertypes, etc.), and the argument and result types of the operations.

1.2 Subtype versus Subclass

It is important to carefully define the terms “type”, “class”, “subtype” and “subclass”. By a type we mean an abstract data type (ADT), as characterized by a (behavioral) specification. By a class we mean an implementation module (such as classes in Smalltalk or C++).

A subclass is formed from another class (its superclass) by inheritance of code; this has no particular relation to behavior, as an object of a subclass may behave quite differently from an object of one of its superclasses. For example the class IntStack may be a subclass of IntDEQueue, although a IntStack object cannot respond to all the messages that one would want to send to a IntDEQueue [Sny86].

If one were to make an analogy between the notions of subtype and subclass, one would say that, by contrast, a subtype is formed from another ADT by inheritance of specifications, not code. For example, one might specify the type IntDEQueue by inheriting from IntStack, even if one chooses to implement the class IntStack as a subclass of the class IntDEQueue. The inheritance analogy is only approximate, however, as subtyping has to do with specified behavior,
not with how that behavior is specified. That is, subtyping is independent of specification inheritance. More precisely, a type \( S \) is a subtype of \( T \) if one can use objects of type \( S \) in a program where objects of type \( T \) are expected without any surprising results. This certainly implies that each object of type \( S \) acts like some object of type \( T \) \cite{SC86} \cite{SCBR86} \cite{Lea89}. In many practical cases it also implies that there is a homomorphic coercion function, \( f_{S,T} \), that maps the values of objects of type \( S \) to values of objects of type \( T \) in such a way that for all instance operations of the supertype and for all \( x \) of the subtype \( S \),

\[
\begin{align*}
\text{SuperPreCond}(f_{S,T}(x)) & \Rightarrow \text{SubPreCond}(x) \\
\text{SuperPostCond}(f_{S,T}(x)) & \Leftarrow \text{SubPostCond}(x)
\end{align*}
\]

where “SuperPreCond” is the precondition of the supertype’s instance operation, etc. \cite{Am87} \cite{Am89} \cite{Am91} \cite{IWR93} \cite{IWR93b}.

### 1.3 Plan

In the following we discuss inheritance of specifications in ISLs. Our ideas come from our work on the ISLs Larch/Smalltalk \cite{Ch91} and Larch/C++ \cite{LC93} \cite{LC93a}, and our work on the semantics of subtyping in OOPLs \cite{Lea89} \cite{LW80} \cite{LW90} \cite{LW90} \cite{LP91} \cite{LW92}.

## 2 Inheritance of Specifications

For an example, consider the types \( \text{BankAccount} \) and \( \text{PlusAccount} \). The supertype, \( \text{BankAccount} \), has just a savings account. The subtype, \( \text{PlusAccount} \), also has a (“free”) checking account. We want to specify instance operations such as \( \text{balance} \) and \( \text{pay_interest} \), for \( \text{BankAccount} \) and have these specifications be inherited by \( \text{PlusAccount} \).

The Larch/C++ interface specification of \( \text{BankAccount} \) is given in Figure 1. The LSL \cite{GHL+93} trait \( \text{BankAccountTrait} \) it uses is presented in Figure 2. The trait \( \text{Rational} \) which is included by \( \text{BankAccountTrait} \), is found in the Larch Shared Language Handbook \cite{GHL+93, ApA}. The member functions are specified as \textit{virtual}, which means that the code executed in a call such as \( \text{ba} \rightarrow \text{pay_interest}() \) will execute code determined by the dynamic class of the object pointed to by \( \text{ba} \).

In order to state the inheritance problem as clearly as possible, we specify the subtype \( \text{PlusAccount} \) in Figure 3 without using inheritance. The LSL trait \( \text{PlusAccountTrait} \) is specified in Figure 4.

### 2.1 The Specification Inheritance Problem

The specification inheritance problem is to state the specification of types like \( \text{PlusAccount} \) as succinctly as possible, and to give the specification with inheri-
class BankAccount {
    uses BankAccountTrait(BankAccount for Acct);
    public:
        BankAccount(double amt) {
            requires (1/100) ≤ rational(amt);
            constructs self;
            ensures approximates(amt, balance(self'), 1/100);
        }
        virtual double balance() const {
            ensures approximates(result, balance(self'), (1/100));
        }
        virtual void pay_interest(double rate) {
            requires (0/1) ≤ rate ∧ rate ≤ (1/1);
            modifies self;
            ensures approximates(toDouble(balance(self')),
                                ((1/1) + rational(rate)) × balance(self'),
                                1/100);
        }
        virtual void update(double amt) {
            modifies self;
            ensures approximates(toDouble(balance(self')),
                                balance(self') + rational(amt), 1/100);
        }
};

Figure 1: Larch/C++ interface specification of the (super)type BankAccount.

BankAccountTrait(Acct): trait
    includes Rational % defines the sort Q
    introduces
        createAcct: Q → Acct
        balance: Acct → Q
    asserts ∀ q: Q
        balance(createAcct(q)) = q

Figure 2: The trait that specifies the abstract values of BankAccount objects.

tance a semantics that matches the intended specification as nearly as possible.
For PlusAccount, what one wants to write is something like the specification in Figure 5. In this specification, the type PlusAccount inherits the specifications of balance, pay_interest, and update. It is hard to imagine the interface specification of PlusAccount being more succinct. The question is, what does
### Possible Semantics of Specification Inheritance

A little reflection is enough to convince one that what should be done is to copy each inherited operation specification from the parent specification to the inheriting type’s specification [CDD+89] [DD90] [Cus91] [LC93b]. The semantic question then becomes: given that the parent’s specification was written in terms of the abstract values of the parent type, how can one interpret it for the abstract values of the inheriting type? We have identified four potential answers to this question in the context of a Larch-style ISL.
PlusAccount Trait: trait
introduces
savNchk: Q,q → PA
checking: PA → Q
savings: PA → Q
asserts
∀ q1,q2: Q, pa: PA
savings(savNchk(q1,q2)) = q1;
checking(savNchk(q1,q2)) = q2

Figure 4: The trait that specifies the abstract values of PlusAccount objects.

class PlusAccount : public BankAccount {
    uses PlusAccount Trait(PlusAccount for PA);
public:
    PlusAccount(double savings_balance, double checking_balance) {
        requires (1/100) ≤ rational(savings_balance)
        ∧ (0/1) ≤ rational(checking_balance);
        constructs self;
        ensures approximates(savings_balance, savings(self'), 1/100);
            ∧ approximates(checking_balance, checking(self'), 1/100);
    }
};

Figure 5: Ideal interface specification of PlusAccount as a subtype of BankAccount, using specification inheritance.

1. Use the same sort of abstract values (i.e., extending the same LSL trait) for the subtype as for the supertype [GM87] [MOM90]. Since the abstract values of the types are the same, there is no problem in interpreting the parent type’s specification.

2. Define a homomorphic1 coercion function that maps the abstract values of the inheriting type to the abstract values of the parent type [Ame87] [Ame89] [Ame91] [LW93a] [LW93b]. The parent type’s specification is interpreted by using this function to coerce the inheriting type’s abstract values to the types assumed in its specification.

1Homomorphic in the sense that it commutes with the trait functions of the parent type; for subtypes it should also commute with the instance operations of the supertype in the sense of Formula (1) and (2) above.
3. Define a homomorphic relation that relates each inheriting type’s abstract value to at least one parent type abstract value, which is used to coerce the abstract values of the inheriting type to the parent type [Lea89]. The parent type’s specification is interpreted by using this relation to obtain a set of parent type abstract values, and these are all used to interpret the parent type’s specification.

4. Overload each trait function that takes an argument of the parent type’s abstract values so that it is defined on abstract values of the inheriting type [LW90] [Lea91] [LW92]. For the abstract values of an inheriting type, these overloaded trait functions are used to interpret the inherited specification.

These approaches are discussed and compared below.

2.2.1 Using the Same Sort of Abstract Values

The approach of using the same sort for the inheriting type’s specification as for the parent type specification is slick when it works. It often works when the inheriting type is a simple restriction on the parent type; for example, when a subtype’s abstract values are a subset of the supertype’s. However, it does not work well when the subtype’s objects contain more information than the supertype’s, as is the case with PlusAccount. While it is always possible to specify both sets of abstract values by using a disjoint union as the abstract value set of the parent type, doing so is not modular.

2.2.2 Using a Coercion Function

In our example, one can define a trait function, toAcct, in PlusAccountTrait that maps values of sort PA to values of sort Acct. There are infinitely many such mappings. The mapping which makes the meaning of Figure 5 closest to the meaning of Figure 3 is the one defined by the following axiom.

\[
toAcct(sawNchck(q1,q2)) = createAcct(q1+q2)
\]

Since there are so many possible mappings, the desired coercion function should be specified for the inheriting type. In Larch/C++ this can be done by specifying toAcct in the trait PlusAccountTrait, and adding to the interface specification of Figure 5 a line of the following form.

\[
simulates BankAccount by toAcct
\]

However the requirement that the coercion be a function, and that it be homomorphic, sometimes makes specification inconvenient. Consider the specification of an abstract class Graph, which has no way to create objects, but is intended as the common supertype of DirectedGraph and UndirectedGraph.
One way to describe the abstract values of \texttt{Graph} is as a pair of sets: a set of nodes and a set of edges. The edges cannot be undirected, as this would make them useless for \texttt{DirectedGraph}. But then one cannot specify the abstract values of \texttt{UndirectedGraph} by identifying the edges \([n,m]\) and \([m,n]\), because doing that makes having a homomorphic function from the abstract values of \texttt{UndirectedGraph} to the abstract values of \texttt{Graph} impossible. So the specifier of \texttt{UndirectedGraph} is forced to specify the abstract values without making this identification, which certainly complicates the specification of \texttt{UndirectedGraph} (see [CL93] for how this is done). Note, however, that this is not a modularity problem.

### 2.2.3 Using a Coercion Relation

A homomorphic relation is a generalization of a homomorphic function. Because it can coerce an inheriting type’s abstract value to a set of abstract values of the supertype (viewing the relation as a set-valued function), one can avoid the inconvenience described in the previous paragraph. However the disadvantage of homomorphic relations is that there is much to prove before one is convinced that assertion evaluation is well-defined, because of the possible ambiguity in dealing with sets of abstract values [Lea89] [Lea90].

### 2.2.4 Overloading the Trait Functions

This approach attacks the problem of how to interpret the parent type’s specification directly. It is clearly more general than the approaches above, because the others also, in effect, overload the trait functions so that they are defined on abstract values of the inheriting type. Permitting the trait functions to be overloaded in any way at all allows more flexibility, which can be exploited to partially solve the information loss problem (described in Section 3 below).

However, we recently discovered that this approach is not completely modular. Consider the specification of a type \texttt{Node}, which is used in the specification of the type \texttt{Graph}. Let us suppose that there is a trait function in the trait defining the abstract values of type \texttt{Graph} called \texttt{includesNode}, with the signature: \texttt{includesNode: Graph, Node \rightarrow Node}. When one defines a subtype of \texttt{Node}, say \texttt{ColoredNode}, then one must overload the trait function \texttt{includesNode} so that it is defined when its second argument is a \texttt{ColoredNode}. But this violates the definition of modularity of specifications, because the specifier of \texttt{ColoredNode} should not have to know about the specification of \texttt{Graph}, which is not a supertype of \texttt{Node}.

Another problem with this approach is that it requires a nonstandard interpretation of equality (=).
2.3 Discussion

Recall that in an OOPL, inheritance of code does not have to be used to make subtypes. Should inheritance of specifications necessarily be used to make subtypes? We can see no good reason for this in general. All of the approaches mentioned above work for inheriting specifications, even if the inheriting type is not a subtype of the parent type. Indeed this is a plus, as one would want the specification to be well-defined so that one can prove whether a claimed subtype relationship is legal or not.

There is also the semantic issue of what to do when multiple specifications of the same operation are inherited from different parent types. For the sake of brevity, we only offer our opinion that the best thing for a specification language to do is to prohibit the use of inheritance for such specifications; after all, the understandability of the specification is not decreased by giving the specification explicitly.

3 The Information Loss Problem

The information loss problem occurs when an inheriting type’s abstract values contain more information than its parent type’s. Consider what the inherited specification of pay interest in Figure 5 says compared to its specification in Figure 3. From Figure 5 one can conclude that the total balance is increased by the specified interest, but the distribution between checkings and savings is not specified. Thus, besides paying interest, an implementation of pay interest for PlusAccount is allowed to transfer money between checking and savings!

For post-conditions of the form self’ = tf(self’), where tf is a trait function, the problem may be solved by specially overloading the trait function tf to avoid the information loss. However, not all post-conditions take this form—witness pay interest. So overloading the trait functions is not a general solution.

Meyer’s OOPL Eiffel has a way to avoid information loss without resorting to complete respecification. In subtypes, the Eiffel specifier can (only) conjoin an additional assertion to the post-condition using the keyword then [Mey92]. For example, one would specify pay interest by specifying that the ratio of the checking to the savings to the checking parts of a PlusAccount is unchanged by pay interest, as in the following.

```c
virtual void pay_interest(double rate) {
    ensures then if checking(self') = 0 then checking(self') = 0
    else (savings(self')/checking(self'))
        = (savings(self')/checking(self'));
}
```

However, there is no reason to limit such shorthands to the specification of subtypes.
The information loss problem may also be amenable to solutions similar to those proposed for the frame problem (see below).

4 The Frame Problem

The frame problem is how to say “and nothing else changes” in a specification [BMR93]. In Larch/C++, a function specification has a modifies clause that says what objects the function is allowed to change. However, when a modifies clause of the form modifies self is inherited, it means that the abstract value as a whole may change. This may be less restrictive than intended, as extra information that in the subtype’s abstract values may or may not be intended to change. The approach advocated in [BMR93] looks promising as a way to solve this.

5 Summary Position

For each inheriting type, the specifier should state a coercion function (or relation). This avoids all modularity problems. For maximum flexibility, a specification language could also allow individual trait functions to be defined, not by the coercion, but by explicit overloading. That way only overloads that are different than those produced by the coercion function need be defined. In effect this gives inheritance with overriding to trait functions.

Allow adding conjuncts to post-conditions to ease the information loss problem without requiring complete respecification.

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