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BLOCK STANLEY DECOMPOSITIONS
I. ELEMENTARY AND GNOMON DECOMPOSITIONS

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ABSTRACT: Stanley decompositions are used in invariant theory and the theory of normal forms for dynamical systems to provide a unique way of writing each invariant as a polynomial in the Hilbert basis elements. Since the required Stanley decompositions can be very long, we introduce a more concise notation called a block decomposition, along with three notions of shortness (incompressibility, minimality of Stanley spaces, and minimality of blocks) for block decompositions. We give two algorithms that generate different block decompositions, which we call elementary and gnomon decompositions, and give examples. Soleyman-Jahan's criterion for a Stanley decomposition to come from a prime filtration is reformulated to apply to block decompositions. We simplify his proof, and apply the theorem to show that elementary and gnomon decompositions come from "subprime" filtrations. In a sequel to this paper we will introduce two additional algorithms that generate block decompositions that may not always be subprime, but are always incompressible.

1. Introduction

Let $I \subset S = K[\mathbf{x}] = K[x_1, \dots, x_n]$ be a monomial ideal and let M be the set of *standard monomials*, that is, the monomials $\mathbf{x}^{\mathbf{k}} = x_1^{k_1} \cdots x_n^{k_n} \in K[\mathbf{x}]$ that are *not* in I . Each coset in S/I has exactly one representative as a *standard polynomial* in $\text{Span}(M)$, and S/I is isomorphic to $\text{Span}(M)$ as a vector space over K , although the ring structure of S/I is lost. A *block decomposition* of M (defined precisely in Section 2) is an expression for M as a disjoint union of rectangular *blocks* of standard monomials, regarded as points in Newton space. A block will be called a *Stanley block* if its span is a Stanley space; a block decomposition of A into Stanley blocks is equivalent to a Stanley decomposition of S/I .

In this paper and its sequel we study algorithms leading from I to a block decomposition for M , emphasizing that different algorithms produce block decompositions with different algebraic and geometric properties. Our long-range goal is to produce block decompositions that are as simple as possible, with respect to various criteria. In this first paper we present two algorithms, producing what we call the *elementary* and *gnomon* block decompositions respectively, and show that these are associated with *subprime filtrations* of I , a natural generalization of prime filtrations. In doing so, we generalize Soleyman-Jahan's criterion in [16, Prop. 2.7] for a Stanley decomposition to be associated with a prime filtration, and simplify the proof. The gnomon decomposition is shorter (contains fewer blocks) than the elementary decomposition. In the second paper we present algorithms leading to what we call the *organized* and *stacked* block decompositions. These are *incompressible*, in the sense that they cannot be shortened by combining blocks. We have not yet found an algorithm that always produces a minimal block decomposition (one with the fewest blocks), but since any such block decomposition must be incompressible, these algorithms are a step in that direction.

From one point of view, block decompositions are a generalization of Stanley decompositions, and are in general *coarser* than Stanley decompositions. From another point of view, each block decomposition has a unique minimal Stanley refinement, given by Algorithm 2.2 below, and since this algorithm is very easy, the block decomposition can be taken as a *compressed notation* for this Stanley decomposition. From this viewpoint, we are studying algorithms leading from I to a Stanley decomposition of S/I . The usual algorithm of this type, [17, Lemma 2.4], proceeds by induction on the variables x_1, \dots, x_n . Our gnomon algorithm works by induction on the generators of I , so that if one has

determined a block decomposition for S/I and then adds a generator to I , it is not necessary to start over from the beginning.

In Section 2 we define block decompositions and Stanley decompositions and present Algorithm 2.2 that relates them. Sections 3 and 4 contain the algorithms for the elementary and gnomon block decompositions. Section 5 contains the generalization of Soleyman-Jahan's theorem and the proof that elementary and gnomon decompositions are subprime.

Stanley decompositions are best known among algebraists in connection with the conjecture that the Stanley depth of a module over a polynomial algebra is an upper bound for the classical depth, but have a separate life in applied mathematics through their use in describing subalgebras. This life is based entirely on the first two pages of [17] and, for one application, on [18]. The motivation for our work comes from that direction. Let \mathcal{A} be a subalgebra of $K[\mathbf{y}] = K[y_1, \dots, y_m]$. It is desired to find a formula, containing finitely many arbitrary polynomial functions, that expresses each element of \mathcal{A} exactly once as the arbitrary functions are varied. A set of polynomials $f_1, \dots, f_n \in \mathcal{A}$ is a *Hilbert basis* for \mathcal{A} if the monomials $f_1^{k_1} \cdots f_n^{k_n}$ span \mathcal{A} as a vector space over K . The expression for a given f in terms of the Hilbert basis is usually not unique, because the monomials $f_1^{k_1} \cdots f_n^{k_n}$ may not be linearly independent. There always exist linearly independent subsets of these monomials, which we call *preferred sets*, and the expression of any f as a linear combination of these is unique. But a preferred set of monomials is infinite, and may not have a finite description leading to an expression involving arbitrary polynomial functions. Stanley decompositions in the manner of [17] provide suitable preferred sets. Let $\Phi : K[\mathbf{x}] \rightarrow K[\mathbf{y}]$ be the unique algebra homomorphism such that $\Phi(x_i) = f_i$; then $J = \ker \Phi$ is the ideal of relations among the f_i . Compute a Gröbner basis for J using an appropriate elimination order on $K[\mathbf{x}, \mathbf{y}]$ ([1, p. 81]), and obtain the monomial ideal I of leading terms of J , which is generated by the leading terms of the Gröbner basis. The standard monomials of I provide a preferred set of monomials in f_1, \dots, f_n when each x_i is replaced by f_i . In fact \mathcal{A} is isomorphic as an algebra to $K[\mathbf{x}]/J$, and isomorphic as a vector space to both $K[\mathbf{x}]/I$ and $\text{Span}(M)$. Furthermore, a Stanley decomposition for $K[\mathbf{x}]/I$ provides the desired formula expressing each element of \mathcal{A} uniquely. For instance, if \mathcal{A} has a Hilbert basis of four elements f_1, \dots, f_4 with the single relation $f_3^2 - f_2^3 - f_1^2 f_4 = 0$, and f_3^2 is taken as the leading term, then $I = \langle x_3^2 \rangle$ and the Stanley decomposition given by the algorithm

of [17] is

$$(1.1) \quad K[x_1, x_2, x_3, x_4]/\langle x_2^3 \rangle \cong K[x_1, x_2, x_4] \oplus K[x_1, x_2, x_4]x_3.$$

It follows that every element of \mathcal{A} can be written uniquely in the form

$$f = F_1(f_1, f_2, f_4) + F_2(f_1, f_2, f_4)f_3,$$

where F_1 and F_2 are arbitrary polynomials in three variables. (This example arises for the classical seminvariants of the binary cubic form.)

Notice that the number of arbitrary polynomials is the same as the number of Stanley spaces. The number of such polynomials can become large in larger problems, and it is natural to ask if this can be reduced. In part this can be done by a good choice of the term order; for instance, if f_2^3 was taken as the leading term in the example, there would be three Stanley spaces instead of two. But even once I is selected, there can be large differences in the number of Stanley terms in different Stanley decompositions for S/I . This suggests the problem: *Find a minimal Stanley decomposition for S/I given a fixed monomial ideal I .* A related, but different, problem is to *find a minimal block decomposition for M given I* ; here the minimality is with regard to the number of blocks, not the number of Stanley spaces.

It is not an accident that the example mentioned above comes from classical invariant theory. Invariant theory is closely linked with normal form theory for systems of differential equations with nilpotent linear part. There was a flurry of work on this topic in the late 1980s and early 1990s ([2], [3], [5], [18], [4], [17]), and a second flurry more recently ([12], [13], [14, Ch.12], [15], [9], [10], [7], [6]). Additional motivation for our questions comes from these papers. The algebra of scalar fields that are invariant under the the group $\{e^{Mt} : t \in \mathbb{R}\}$, where M is a nilpotent matrix, plays an important role. These invariants, in turn, are the same as classical seminvariants of a binary form, or joint seminvariants of several forms. Knowing the seminvariants of several forms, *transvectants* are used to find a Hilbert basis for the the joint seminvariants. The *box product* method introduced in [13] lifts this procedure to the level of Stanley decompositions. The box product of two Stanley decompositions for the seminvariants of two binary forms is a Stanley decomposition for the joint seminvariants of the two forms; thus the uniqueness issue is taken care of automatically, without further Gröbner basis work. In [?] I (J.M.) have reformulated the box product as a product of block decompositions, in which \boxtimes distributes over disjoint union; the basic unit of computation is the box product of two blocks, for which algorithms are given, and the final result (expressed in Stanley decompositions) is shorter than for the method in [13]. It is

best, then, to start with block decompositions with a minimal number of blocks.

This research began when I (J.M.) asked T.M., a software engineer, to write programs for me concerning Stanley decompositions. He created block decompositions as a computer-friendly shorthand for Stanley decompositions, and the gnomon algorithm to provide an interactive program in which ideal generators could be added one at a time. I provided the mathematical write-up, proofs, and examples, and the material on subprime filtrations. (The word *subprime* was provided by the mortgage crisis.)

2. Block decompositions

Let $V \subset S = K[\mathbf{x}]$ be a vector subspace; V is a **monomial space** if it is spanned by the set $M = \text{Mon}(V)$ of monomials in V , which is then also the unique monomial basis for V . Under the one-to-one correspondence $V = \text{Span}(M)$, $M = \text{Mon}(V)$ between monomial spaces and their monomial bases, a direct sum of two spaces corresponds to the disjoint union of their bases. We often identify $\mathbf{x}^{\mathbf{m}}$ with $\mathbf{m} = (m_1, \dots, m_n)$ in the **Newton space** $\mathbb{Z}_{\geq 0}^n$.

A **Stanley space** is a monomial space of the form

$$K[X]\varphi = \{f\varphi : f \in K[X]\},$$

where X is a subset of the set of variables $\{x_1, \dots, x_n\}$ and φ is a monomial. An expression exhibiting a monomial space V as a direct sum of finitely many Stanley spaces is called a **Stanley decomposition** of V . (Because of this finiteness requirement, there are monomial spaces, such as $\text{Span}\{1, x^2, x^4, x^6, \dots\} \subset K[x]$, that do not have Stanley decompositions.)

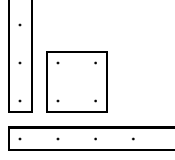
Let I be a monomial ideal in S and let $M = \text{Mon}(S) \setminus \text{Mon}(I)$ be the set of **standard monomials** for I . Then $V = \text{Span}(M)$ is the space of **standard polynomials**, and as noted in the introduction, $S/I \cong V$ as vector spaces. If $S = K[x, y]$ and $I = \langle xy^3, x^3y \rangle$, a Stanley decomposition for S/I is

$$(2.1) \quad K[x] \oplus K[y]y \oplus Kxy \oplus Kx^2y \oplus Kxy^2 \oplus Kx^2y^2.$$

Applying the correspondences above, we replace each Stanley space by its set of monomials and replace \oplus by the disjoint union symbol \sqcup , obtaining

$$\{1, x, x^2, \dots\} \sqcup \{y, y^2, y^3, \dots\} \sqcup \{xy, xy^2, x^2y, x^2y^2\}.$$

In this way the Stanley decomposition (2.1) becomes the following *block decomposition* in Newton space:



More generally, let $[a, b]$ denote an integer interval, with the understanding that $[a, \infty]$ does not include ∞ and that $[a, b] = \emptyset$ if $b < a$. Define a **block** B in $\mathbb{Z}_{\geq 0}^n$ to be a Cartesian product of n intervals:

$$(2.2) \quad B = \begin{bmatrix} \mathbf{b} \\ \mathbf{a} \end{bmatrix} = \begin{bmatrix} b_1 & b_2 & \cdots & b_n \\ a_1 & a_2 & \cdots & a_n \end{bmatrix} = [a_1, b_1] \times \cdots \times [a_n, b_n].$$

If any $b_i < a_i$, the block is empty. The bottom row \mathbf{a} of a block is called its **inner corner**. The letter B will be used ambiguously to mean the matrix as such, the set of integer vectors represented by the matrix, and the set of monomials having those integer vectors as exponents. Let $M \subset \text{Mon}(S)$. A **block decomposition** of M is an expression

$$(2.3) \quad M = B^1 \sqcup B^2 \sqcup \cdots \sqcup B^s,$$

where each B^k is a block. For example, the Stanley decomposition (2.1) can be written as the block decomposition

$$(2.4) \quad B^1 \sqcup B^2 \sqcup B^3 = \begin{bmatrix} \infty & 0 \\ 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 0 & \infty \\ 0 & 1 \end{bmatrix} \sqcup \begin{bmatrix} 2 & 2 \\ 1 & 1 \end{bmatrix},$$

corresponding to the diagram shown above in Newton space.

2.1. Remark. After this research was mostly completed, we learned of the **interval decompositions** defined in [8]. They define the **interval** $[\mathbf{a}, \mathbf{b}]$ to be the set of monomials lying between $\mathbf{x}^{\mathbf{a}}$ and $\mathbf{x}^{\mathbf{b}}$ in division partial order. This is essentially the same as our block B in (2.4), except that they consider only monomials that divide a chosen $\mathbf{x}^{\mathbf{c}}$, so that, in particular, ∞ is not allowed in the top row. This is sufficient since they only consider interval decompositions for I/J , where I and J are two monomial ideals in S .

A block decomposition is called **compressible** if the union of some subset of the blocks is itself a block. In this case the decomposition can be simplified by performing the union. However, performing one such union may prevent another one, so a given decomposition may be compressible in several ways to give distinct incompressible decompositions. A decomposition is called **minimal** if there is no decomposition (describing the same set of monomials) having fewer blocks. Incompressibility and minimality are two possible notions of “simplicity” of a decomposition (as briefly discussed in the introduction). An incompressible decomposition need not be minimal, and a minimal decomposition need not be unique. The following example shows that

incompressibility is a global property of a decomposition, and cannot be detected by examining the decomposition locally: The disjoint union

$$\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 0 & 2 \\ 0 & 1 \end{bmatrix} \sqcup \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} \sqcup \begin{bmatrix} 2 & 1 \\ 2 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 2 & 2 \\ 1 & 2 \end{bmatrix}$$

is compressible to the single block $\begin{bmatrix} 2 & 2 \\ 0 & 0 \end{bmatrix}$, although no proper subset of the five blocks can be compressed.

A **Stanley block** is a block (2.2) in which each b_i equals either a_i or ∞ . The span of a Stanley block is a Stanley space, and any Stanley decomposition can be changed to a block decomposition by replacing Stanley spaces by Stanley blocks and \oplus by \sqcup . Such a block decomposition can almost always be compressed. For example, $K[x_1, x_2, x_4] \oplus K[x_1, x_2, x_4]x_3$ becomes

$$\begin{bmatrix} \infty & \infty & 0 & \infty \\ 0 & 0 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & \infty & 1 & \infty \\ 0 & 0 & 1 & 0 \end{bmatrix} = \begin{bmatrix} \infty & \infty & 1 & \infty \\ 0 & 0 & 0 & 0 \end{bmatrix}.$$

Any block B can be written as a disjoint union of its **bounded part** and its **unbounded part**; the bounded part is obtained by replacing the unbounded columns with zeroes, the unbounded part by doing the same with the bounded columns. For instance,

$$B = \begin{bmatrix} \infty & 5 & \infty & 2 \\ 1 & 3 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 0 & 5 & 0 & 2 \\ 0 & 3 & 0 & 1 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 0 & \infty & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}.$$

Let X be the set of variables associated with unbounded columns of B (x_1 and x_3 in the example). Let ψ be the monomial in the bottom row of the unbounded part ($x_2^2x_4$ in the example). Let $\theta_1, \dots, \theta_s$ be the finite set of monomials in the bounded part ($x_2^3x_4, x_2^4x_4, x_2^5x_4, x_2^3x_4^2, x_2^4x_4^2, x_2^5x_4^2$, with $s = 6$). Let $\varphi_k = \theta_k\psi$ for $k = 1, \dots, s$. Then the block B is converted to its minimal Stanley decomposition as follows.

2.2. Algorithm. *The Stanley decomposition with the fewest Stanley spaces that describes the span of the monomials in B is*

$$\text{Span}(B) = K(X)\varphi_1 \oplus \dots \oplus K(X)\varphi_s.$$

Note that all of the coefficient algebras $K(X)$ in the minimal Stanley decomposition of a block are equal, and that the number of Stanley spaces s is the product of the numbers of elements in the intervals defined by the bounded columns of B . When Algorithm 2.2 is applied to each block in a block decomposition, the coefficient rings may differ from block to block, but the total number of Stanley spaces is just the sum of the number for each block.

As an introduction to the next two sections, we now give examples of the *elementary* and *gnomon* block decompositions defined in those sections. For

$$(2.5) \quad I = \langle x^3y^9, x^7y^5 \rangle \subset S = K[x, y],$$

the *elementary* decomposition of S/I is the disjoint union of the blocks

$$(2.6) \quad \begin{bmatrix} 2 & 4 \\ 0 & 0 \end{bmatrix}, \begin{bmatrix} 2 & 8 \\ 0 & 5 \end{bmatrix}, \begin{bmatrix} 2 & \infty \\ 0 & 9 \end{bmatrix}, \begin{bmatrix} 6 & 4 \\ 3 & 0 \end{bmatrix}, \begin{bmatrix} 6 & 8 \\ 3 & 5 \end{bmatrix}, \begin{bmatrix} \infty & 4 \\ 7 & 0 \end{bmatrix}.$$

The *gnomon* decomposition comes in two forms,

$$(2.7) \quad \begin{bmatrix} 2 & \infty \\ 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 6 & 8 \\ 3 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 4 \\ 7 & 0 \end{bmatrix}$$

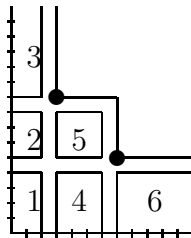
and

$$(2.8) \quad \begin{bmatrix} \infty & 4 \\ 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 6 & \infty \\ 0 & 5 \end{bmatrix} \sqcup \begin{bmatrix} 2 & \infty \\ 0 & 9 \end{bmatrix}.$$

The elementary decomposition is compressible to either of the gnomon decompositions, and (in this case, but not always) the gnomon decompositions are incompressible.

3. THE ELEMENTARY BLOCK DECOMPOSITION

In Newton space, the **division partial order** $\mathbf{x}^{\mathbf{m}} | \mathbf{x}^{\mathbf{m}'}$ will be written as $\mathbf{m} \preceq \mathbf{m}'$, meaning $m_i \leq m'_i$ for $i = 1, \dots, n$. For variables in Newton space we use $\boldsymbol{\mu} = (\mu_1, \dots, \mu_n)$, reserving $\mathbf{m} = (m_1, \dots, m_n)$ for constants, so that an equation such as $\mu_2 = m_2$ will represent the “hyperplane” through \mathbf{m} perpendicular to the μ_2 axis in $\mathbb{Z}_{\geq 0}^n$. Let $S = K[\mathbf{x}]$ and let $I = \langle \mathbf{x}^{\mathbf{m}^1}, \dots, \mathbf{x}^{\mathbf{m}^r} \rangle \subset S$ be the monomial ideal with the indicated generators; we also write $I = \langle \mathbf{m}^1, \dots, \mathbf{m}^r \rangle$. It is assumed that these generators are minimal, that is, no redundant generators (divisible by another generator) are included. The **elementary block decomposition** for S/I is created by first gridding the Newton space $\mathbb{Z}_{\geq 0}^n$ with the hyperplanes $\mu_i = m_i^r$ for $i = 1, \dots, n$ and $j = 1, \dots, s$, passing through all points of the minimal generating set, and then discarding those of the resulting blocks that belong to $\text{Mon}(I)$ rather than $\text{Mon}(S) \setminus \text{Mon}(I)$. The following figure illustrates the elementary decomposition (2.6) at the end of Section 2. The solid dots are the generators at $(3, 9)$ and $(7, 5)$, the dark lines are the boundary of $\text{Mon}(I)$, and the blocks are numbered in the order that they appear in (2.6).



The elementary block decomposition is almost always compressible, and so is not very desirable in itself, but it shows that a natural block decomposition always exists. The other decompositions (gnomon, organized, and stacked) in this paper and its sequel can be obtained by compressing the elementary decomposition, so we will never need to consider blocks smaller than the blocks in the elementary decomposition, or blocks that are not disjoint unions of elementary blocks.

The following is a precise algorithm to create the elementary decomposition.

3.1. Algorithm. Let $I = \langle \mathbf{m}^1, \dots, \mathbf{m}^r \rangle$ be presented by its minimal generating set. The elementary block decomposition is generated as follows.

1. For $i = 1, \dots, n$, create the list of exponents of x_i in the set of generators, adding zero at the beginning of each list:

$$L_1 : 0, m_1^1, \dots, m_1^r;$$

$$L_2 : 0, m_2^1, \dots, m_2^r;$$

⋮

$$L_n : 0, m_n^1, \dots, m_n^r.$$

2. Delete any repetitions in each list L_i , and re-order each list in increasing numerical order.
3. Create a preliminary list of inner corners as follows. For $i = 1, \dots, n$, choose an entry a_i from L_i ; then create an inner corner $\mathbf{a} = (a_1, \dots, a_n)$. Do this in all possible ways.
4. Refine the list of inner corners by discarding any $\mathbf{a} \in I$, that is, any \mathbf{a} such that $\mathbf{m}^i \prec \mathbf{a}$ for some i .
5. Create a block B for each inner corner \mathbf{a} in the refined list, by determining its outer corner \mathbf{b} , as follows. If a_i has a successor a'_i in L_i (under the ordering from step 2), put $b_i = a'_i - 1$. If a_i is the last element of L_i , but $b_i = \infty$.
6. Let s be the number of resulting blocks, and enumerate these as

$$B^k = \begin{bmatrix} \mathbf{b}^k \\ \mathbf{a}^k \end{bmatrix}$$

for $k = 1, \dots, s$.

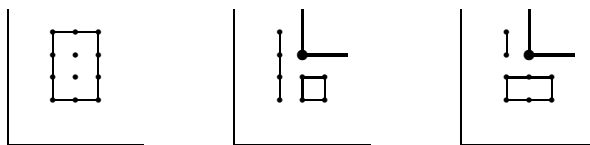
3.2. Remark. The order of enumeration in Step 6 does not matter except in Section 5. For that section, enumerate the blocks in lexicographic order by their inner corners.

For the example (2.5), the lists from Step 2 are $L_1 = \{0, 3, 7\}$, $L_2 = \{0, 5, 9\}$. The preliminary list of inner corners (in lexicographic order) is $\{(0, 0), (0, 5), (0, 9), (3, 0), (3, 5), (3, 9), (7, 0), (7, 5), (7, 9)\}$. From this we discard $(3, 9)$, $(7, 5)$, and $(7, 9)$. It is now easy to check that the outer corners in (2.6) follow from step 5 above.

3.3. Remark. The elementary decomposition has the following property, which will be used in Section 5. Let $\mathbf{c} \notin I$ be a monomial. Then there exists a unique inner corner \mathbf{a}^k of the elementary decomposition for S/I that is maximal (under \prec) among all inner corners satisfying $\mathbf{a}^k \preceq \mathbf{c}$. The inner corner \mathbf{a}^k satisfying this condition determines the block B^k that contains \mathbf{c} . (The proof is trivial: The block containing \mathbf{c} has for the i -th component of its inner corner the largest element of L_i that is $\leq c_i$.)

4. THE GNOMON DECOMPOSITION

Consider the following **Block Subtraction Problem**: Given a block (2.2) and a principal monomial ideal $I = \langle \mathbf{m} \rangle$, find a block decomposition of $B \setminus I$. If \mathbf{m} is within the block B , there are $n!$ distinct natural solutions. For $n = 2$, the following figure shows a block B in $\mathbb{Z}_{\geq 0}^2$, an ideal I generated by the heavy dot, and the two block decompositions of $B \setminus I$. The set of points in $B \setminus I$ is the sort of L-shaped region the Pythagoreans called a **gnomon** ($\gamma\nu\omega\mu\omega\nu$).



The first of these decompositions, expressed algebraically, takes the form

$$(4.1) \quad \begin{bmatrix} b_1 & b_2 \\ a_1 & a_2 \end{bmatrix} \setminus \langle (m_1, m_2) \rangle = \begin{bmatrix} m_1 - 1 & b_2 \\ a_1 & a_2 \end{bmatrix} \sqcup \begin{bmatrix} b_1 & m_2 - 1 \\ m_1 & a_2 \end{bmatrix}.$$

We will generalize this formula to arbitrary n in such a way that it gives all $n!$ solutions, one for each permutation of the variables x_1, \dots, x_n . First we focus on the natural order of variables (the trivial permutation), which for $n = 2$ gives (4.1).

The following solution for $n = 2$ is valid whether or not \mathbf{m} is in the interior of B , and reduces to (4.1) when it is:

$$\begin{bmatrix} b_1 & b_2 \\ a_1 & a_2 \end{bmatrix} \setminus \langle (m_1, m_2) \rangle = \begin{bmatrix} \min\{b_1, m_1 - 1\} & b_2 \\ a_1 & a_2 \end{bmatrix} \sqcup \begin{bmatrix} b_1 & \min\{b_2, m_2 - 1\} \\ \max\{a_1, m_1\} & a_2 \end{bmatrix}.$$

The first block consists of all monomials in B that are not divisible by $x_1^{m_1}$, and is empty if $m_1 \leq a_1$. The second block consists of those monomials in B that *are* divisible by $x_1^{m_1}$ but not by $x_2^{m_2}$. But this solution has a defect: it sometimes divides B unnecessarily into two nonempty blocks whose union is again B . This happens whenever $a_1 < m_1 < b_1$ but $b_2 < m_2$. In fact, if $m_i > b_i$ for $i = 1$ or 2 or both, then $B \cap I = \emptyset$ and $B \setminus I = B$. This observation also simplifies the formula in the remaining cases (where $m_i \leq b_i$ for all i), because then $\min\{b_i, m_i - 1\}$ can be replaced by $m_i - 1$.

Generalizing to higher dimensions, we define the **gnomon** determined by B and I to be the set represented by the $3 \times n$ matrix

$$\begin{bmatrix} \mathbf{b} \\ \mathbf{m} \\ \mathbf{a} \end{bmatrix} = \begin{bmatrix} b_1 & \cdots & b_n \\ m_1 & \cdots & m_n \\ a_1 & \cdots & a_n \end{bmatrix}$$

as follows. (Placing \mathbf{m} between \mathbf{a} and \mathbf{b} is intended to suggest “cutting B by \mathbf{m} ”.)

1. If there exists $i \in \{1, \dots, n\}$ such that $b_i < m_i$ (that is, if $\mathbf{m} \not\preceq \mathbf{b}$), then

$$(4.2) \quad \begin{bmatrix} \mathbf{b} \\ \mathbf{m} \\ \mathbf{a} \end{bmatrix} = \begin{bmatrix} \mathbf{b} \\ \mathbf{a} \end{bmatrix}.$$

2. If for all i , $m_i \leq b_i$ ($\mathbf{m} \preceq \mathbf{b}$), let

$$(4.3) \quad c_i = \max\{a_i, m_i\}.$$

Then

$$(4.4) \quad \begin{bmatrix} \mathbf{b} \\ \mathbf{m} \\ \mathbf{a} \end{bmatrix} = \begin{bmatrix} m_1 - 1 & b_2 & \cdots & b_n \\ a_1 & a_2 & \cdots & a_n \end{bmatrix} \\ \sqcup \begin{bmatrix} b_1 & m_2 - 1 & b_3 & \cdots & b_n \\ c_1 & a_2 & a_3 & \cdots & a_n \end{bmatrix} \\ \sqcup \begin{bmatrix} b_1 & b_2 & m_3 - 1 & b_4 & \cdots & b_n \\ c_1 & c_2 & a_3 & a_4 & \cdots & a_n \end{bmatrix} \\ \sqcup \cdots \sqcup \begin{bmatrix} b_1 & b_2 & \cdots & b_{n-1} & m_n - 1 \\ c_1 & c_2 & \cdots & c_n & a_n \end{bmatrix}.$$

The following lemma shows that the solution to the Block Subtraction Problem is a gnomon as defined above.

4.1. **Lemma.** *For all nonnegative values of a_i , b_i , and m_i ,*

$$(4.5) \quad \begin{bmatrix} \mathbf{b} \\ \mathbf{a} \end{bmatrix} \setminus \langle \mathbf{m} \rangle = \begin{bmatrix} \mathbf{b} \\ \mathbf{m} \\ \mathbf{a} \end{bmatrix}.$$

Proof. Assume first that $b_i < m_i$ for some i . Then the ideal $\langle \mathbf{m} \rangle$ does not intersect the given block, and (4.5) follows from (4.2).

Next assume that $a_i < m_i \leq b_i$ for all i , so that $c_i = m_i$. Then the original block can be decomposed into two subblocks by the hyperplane $\mu_1 = m_1$:

$$\begin{bmatrix} \mathbf{b} \\ \mathbf{a} \end{bmatrix} = \begin{bmatrix} m_1 - 1 & b_2 & \cdots & b_n \\ a_1 & a_2 & \cdots & a_n \end{bmatrix} \sqcup \begin{bmatrix} b_1 & b_2 & b_3 & \cdots & b_n \\ c_1 & a_2 & a_3 & \cdots & a_n \end{bmatrix}.$$

(Note that c_1 occurs in the lower left corner of the second matrix where m_1 is expected, but these are equal.) The first of these subblocks lies outside of $\langle \mathbf{x}^{\mathbf{m}} \rangle$, and belongs to the difference we are computing; it equals the first block of (4.4). The second subblock is carried forward to the next step in a recursive process: It is split by the hyperplane $\mu_2 = m_2$ into two subblocks, one having $a_2 \leq \mu_2 < m_2$ and the other having $m_2 \leq \mu_2 \leq b_2$. Again, the first subblock is retained and the second carried forward to the third step, and so forth.

Finally, if $m_1 \leq a_1$, then $m_1 - 1 < a_1$ and the first subblock in (4.6) is empty. Furthermore the lower left entry in the second matrix should be a_1 (rather than m_1 , which would cause the block to include monomials that are not in the original block). But in this case $c_1 = a_1$, so (4.6) remains correct. The same reasoning applies if the condition $m_i \leq a_i$ is encountered later in the recursion. \square

Now let $I = \langle \mathbf{m}^1, \dots, \mathbf{m}^r \rangle$ be a monomial ideal for which a block decomposition $B^1 \sqcup \dots \sqcup B^s$ is known, and let $\langle \mathbf{m} \rangle$ be a principal monomial ideal. Let $I' = I + \langle \mathbf{m} \rangle$ be the sum of these ideals (in the usual sense), and note that $I' = \langle \mathbf{m}^1, \dots, \mathbf{m}^r, \mathbf{m} \rangle$. We refer to this as **adding a generator** to I .

4.2. Algorithm. *A block decomposition for S/I' is given by*

$$(4.6) \quad (B^1 \setminus \langle \mathbf{m} \rangle) \sqcup \dots \sqcup (B^s \setminus \langle \mathbf{m} \rangle).$$

Proof. The standard monomials of I' are the monomials that are standard for I and, in addition, are not divisible by $\mathbf{x}^{\mathbf{m}}$. We remove the monomials divisible by $\mathbf{x}^{\mathbf{m}}$ from each block B^k of standard monomials for I by finding the gnomons $B^k \setminus \langle \mathbf{m} \rangle$. Since these are disjoint, their union is a block decomposition. \square

4.3. Remark. For Section 5, where the order of terms in a decomposition matters, we write the result of Algorithm 4.2 in the form

$$\bigsqcup_{k=1}^s \left(\bigsqcup_{i=1}^n B^{ki} \right) = (B^{11} \sqcup \dots \sqcup B^{1n}) \sqcup \dots \sqcup (B^{s1} \sqcup \dots \sqcup B^{sn}).$$

Here B^{k1}, \dots, B^{kn} are the terms of $B^k \setminus \langle \mathbf{m} \rangle$ in the order shown in equation (4.4), retaining any empty blocks, or, if (4.2) applies, then $B^{k1} = B^k$ and $B^{ki} = \emptyset$ for $i = 2, \dots, n$. Retaining the empty sets allows a uniform notation in the proof of Theorem 5.5.

Let $I = \langle \mathbf{m}^1, \dots, \mathbf{m}^r \rangle$ be a monomial ideal given by its minimal generators. The **gnomon decomposition** of S/I , with respect to the standard order x_1, \dots, x_n of the variables, is defined to be the result of applying Algorithm (4.4) below. The result may depend on the order in which the generators are indexed; this **noncommutativity of the generators** is illustrated in the examples below.

4.4. Algorithm. *Apply Algorithm 4.2 repeatedly, beginning with $I = \emptyset$ and adding the generators $\mathbf{x}^{\mathbf{m}^1}, \dots, \mathbf{x}^{\mathbf{m}^r}$ in that order.*

Let π be a permutation of the integers $(1, \dots, n)$, regarded as specifying an order $x_{\pi_1}, \dots, x_{\pi_n}$ of the variables. The **π -gnomon decomposition** of S/I is defined by changing variables temporarily to $y_i = x_{\pi_i}$, applying algorithm (4.4) using the new variables, and then returning to the original variables. (The generators are still taken in their indicated order.)

The blocks of the gnomon decomposition are unions of blocks of the elementary decomposition, because whenever a block is subdivided in Algorithm 4.4, the subdivision occurs along a portion of a hyperplane

$\mu_i = m_i^j$ for some i and j . The subdivisions for elementary blocks occur along these same hyperplanes (but along the entire hyperplane, not a portion of it.)

4.5. Example. For the example (2.5), $I = \langle (3, 9), (7, 5) \rangle$. According to Algorithm 4.4, we begin with the decomposition $\begin{bmatrix} \infty & \infty \\ 0 & 0 \end{bmatrix}$ for $K[x, y]$ and add the generator $(3, 9)$ using Algorithm 4.2, which in turn calls for using Lemma 4.1 to compute

$$\begin{bmatrix} \infty & \infty \\ 0 & 0 \end{bmatrix} \setminus \langle (3, 9) \rangle = \begin{bmatrix} 2 & \infty \\ 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 8 \\ 3 & 0 \end{bmatrix}.$$

The first block contains the monomials that are not divisible by x^3 , the second those that are divisible by x^3 but not by y^9 . Next we find

$$\begin{bmatrix} 2 & \infty \\ 0 & 0 \end{bmatrix} \setminus \langle (7, 5) \rangle = \begin{bmatrix} 2 & \infty \\ 0 & 0 \end{bmatrix}$$

since $2 < 7$ and $(7, 5)$ is not in the block. Finally

$$\begin{bmatrix} \infty & 8 \\ 3 & 0 \end{bmatrix} \setminus \langle x^7 y^5 \rangle = \begin{bmatrix} 6 & 8 \\ 3 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 4 \\ 7 & 0 \end{bmatrix}.$$

The result of these calculations is (2.7).

The same problem works out as follows if we take $I = \langle (7, 5), (3, 9) \rangle$, that is, add the generators in the opposite order:

$$\begin{aligned} \begin{bmatrix} \infty & \infty \\ 7 & 5 \\ 0 & 0 \end{bmatrix} &= \begin{bmatrix} 6 & \infty \\ 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 4 \\ 7 & 0 \end{bmatrix} \\ \begin{bmatrix} 6 & \infty \\ 3 & 9 \\ 0 & 0 \end{bmatrix} &= \begin{bmatrix} 2 & \infty \\ 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 6 & 8 \\ 3 & 0 \end{bmatrix} \\ \begin{bmatrix} \infty & 4 \\ 3 & 9 \\ 7 & 0 \end{bmatrix} &= \begin{bmatrix} \infty & 4 \\ 7 & 0 \end{bmatrix}. \end{aligned}$$

The last line is an instance of equation (4.2). The result is again (2.7), so there is no noncommutativity in the generators in this example. The π -gnomon decomposition with $\pi = (2, 1)$ is (2.8).

4.6. Example. To show noncommutativity of generators, the gnomon decomposition of $K[x, y, z]/\langle (5, 3, 7), (10, 6, 2) \rangle$ is

$$\begin{bmatrix} 4 & \infty & \infty \\ 0 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 2 & \infty \\ 5 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 9 & \infty & 6 \\ 5 & 3 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 5 & 6 \\ 10 & 3 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & \infty & 1 \\ 10 & 6 & 0 \end{bmatrix},$$

which is incompressible; the opposite order of generators gives instead

$$\begin{aligned} & \begin{bmatrix} 4 & \infty & \infty \\ 0 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 9 & 2 & \infty \\ 5 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} 9 & \infty & 6 \\ 5 & 3 & 0 \end{bmatrix} \\ & \sqcup \begin{bmatrix} \infty & 2 & \infty \\ 10 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 5 & 6 \\ 10 & 3 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & \infty & 1 \\ 10 & 6 & 0 \end{bmatrix}. \end{aligned}$$

This is compressible (to the previous result), since

$$\begin{bmatrix} 9 & 2 & \infty \\ 5 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 2 & \infty \\ 10 & 0 & 0 \end{bmatrix} = \begin{bmatrix} \infty & 2 & \infty \\ 5 & 0 & 0 \end{bmatrix}.$$

4.7. Example. The next example will be used in Remark 5.6 to illustrate Theorem 5.5. Let $S = K[x, y, z]$ and $I = \langle z^5, y^2z^3, x^3yz \rangle$. Then

$$\begin{aligned} & \begin{bmatrix} \infty & \infty & \infty \\ 0 & 0 & 5 \\ 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} \infty & \infty & 4 \\ 0 & 0 & 0 \end{bmatrix} \\ & \begin{bmatrix} \infty & \infty & 4 \\ 0 & 2 & 3 \\ 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} \infty & 1 & 4 \\ 0 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & \infty & 2 \\ 0 & 2 & 0 \end{bmatrix} \\ & \begin{bmatrix} \infty & 1 & 4 \\ 3 & 1 & 1 \\ 0 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 2 & 1 & 4 \\ 0 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 0 & 4 \\ 3 & 0 & 0 \end{bmatrix} \sqcup \begin{bmatrix} \infty & 1 & 0 \\ 3 & 1 & 0 \end{bmatrix} \\ & \begin{bmatrix} \infty & \infty & 2 \\ 3 & 1 & 1 \\ 0 & 2 & 0 \end{bmatrix} = \begin{bmatrix} 2 & \infty & 2 \\ 0 & 2 & 0 \end{bmatrix} \sqcup \emptyset \sqcup \begin{bmatrix} \infty & \infty & 0 \\ 3 & 2 & 0 \end{bmatrix}. \end{aligned}$$

The gnomon decomposition for S/I is the union of the last two lines.

We conclude this section with two lemmas that will be used in Section 5.

4.8. Lemma. *With c_i as in (4.3),*

$$(4.7) \quad \begin{bmatrix} b_1 & \cdots & b_n \\ a_1 & \cdots & a_n \end{bmatrix} \cap \langle x_1^{m_1} \cdots x_n^{m_n} \rangle = \begin{bmatrix} b_1 & \cdots & b_n \\ c_1 & \cdots & c_n \end{bmatrix}.$$

Proof. If $m_i > b_i$ for any i , both sides are empty. Otherwise, continue the reasoning in the proof of Lemma 4.1. Observe that the right-hand side of 4.7 would be the next matrix added to (4.4) if the pattern of (4.4) is continued. \square

4.9. Lemma. *Assume that $\mathbf{m} \preceq \mathbf{b}$, so that (4.4) is valid. Let the blocks on the right-hand side of (4.4) be labeled B^1, \dots, B^n . Then for*

any $i = 1, \dots, n$ we have

$$B^i \sqcup \dots \sqcup B^n \sqcup \begin{bmatrix} \mathbf{b} \\ \mathbf{c} \end{bmatrix} = \begin{bmatrix} b_1 & \dots & b_{i-1} & b_i & \dots & b_n \\ c_1 & \dots & c_{i-1} & a_i & \dots & a_n \end{bmatrix}.$$

Proof. Compute the unions successively starting from the right and working to the left, noticing that each union is between matrices that differ only in one column. For instance, if $n = 3$ the first step is

$$\begin{bmatrix} b_1 & b_2 & m_3 - 1 \\ c_1 & c_2 & a_3 \end{bmatrix} \sqcup \begin{bmatrix} b_1 & b_2 & b_3 \\ c_1 & c_2 & c_3 \end{bmatrix} = \begin{bmatrix} b_1 & b_2 & b_3 \\ c_1 & c_2 & a_3 \end{bmatrix}.$$

If $c_3 = m_3$ the result is clear; if $c_3 = a_3$ then the left-hand block is empty and again the result is clear. (The lemma and proof remain valid without assuming $\mathbf{m} \preceq \mathbf{b}$, but in that case it does not apply to the gnomon decomposition, which is governed instead by (4.2).) \square

5. SUBPRIME FILTRATIONS

A **filtration** of S by monomial ideals is a finite nested sequence

$$(5.1) \quad I = I_0 \subset I_1 \subset \dots \subset I_s = S.$$

This is often called a filtration of S/I , and it does induce an actual filtration of S/I by subrings, namely $\{0\} \subset I_1/I \subset I_2/I \subset \dots \subset S/I$. If each set difference $\text{Mon}(I_k) \setminus \text{Mon}(I_{k-1})$ is a block B_k , we call (5.1) a **subprime filtration**. If, further, each block B_k is a Stanley block, we call (5.1) a **prime filtration**; this is simpler than the usual definition in the literature, but is equivalent, as pointed out at the end of this section. A block decomposition, *with the blocks indexed in a specific order*, will be called **subprime** (or **prime**) if it is associated in this way with a subprime (or prime) filtration. In this section all block decompositions are taken as **ordered block decompositions**. This simplifies the discussion, in that we need not consider all possible orderings before deciding whether a given block decomposition is subprime.

The following equations hold trivially for subprime filtrations. Quotients, direct sums, and isomorphisms are to be understood as quotients, direct sums, and isomorphisms of vector spaces, and M is the

set $\text{Mon}(S) \setminus \text{Mon}(I)$ of standard monomials for I in S .

$$(5.2) \quad I_k/I_{k-1} \cong \text{Span}B_k$$

$$(5.3) \quad S/I \cong \text{Span}B_1 \oplus \cdots \oplus \text{Span}B_s$$

$$(5.4) \quad M = B_1 \sqcup \cdots \sqcup B_s$$

$$(5.5) \quad \text{Mon}(I_{k-1}) \sqcup B_k = \text{Mon}(I_k)$$

$$(5.6) \quad I_{k-1} \oplus \text{Span}(B_k) = I_k.$$

That is, with each subprime filtration of S/I is associated a block decomposition in which the blocks are naturally ordered so that the i -th block $B_i = \text{Mon}(I_i) \setminus \text{Mon}(I_{i-1})$, and when the block spaces are successively added (by direct sum) to the I (considered as a vector space), the ideals of the filtration are reconstructed. Thus the subprime filtration and the (ordered) block decomposition contain the same information. Prime filtrations and (ordered) Stanley decompositions are related in the same way.

A convenient notation to exhibit the relationship between ideals and blocks in a subprime filtration is

$$\begin{array}{ccccccccccc} I = I_0 & \subset & I_1 & \subset & I_2 & \cdots & I_{s-1} & \subset & I_s = S. \\ & & \downarrow & & \downarrow & & & & \downarrow \\ & & B_1 & & B_2 & & & & B_s \end{array}$$

For example, the block decomposition (2.4) is associated with a subprime filtration as follows:

$$\begin{array}{ccccccc} \langle xy^3, x^3y \rangle & \subset & \langle xy \rangle & \subset & \langle y \rangle & \subset & \langle 1 \rangle, \\ & & \downarrow & & \downarrow & & \downarrow \\ & & B_1 & & B_2 & & B_3 \end{array}$$

with $B_1 = B^3$, $B_2 = B^2$, and $B_3 = B^1$.

5.1. Remark. Observe that the ideals in the filtration can be found by adding the inner corner of each block to the generators of the previous ideal. For instance, adding the inner corner $(1, 1)$ of $B_1 = B^3$ to $\langle xy^3, x^3y \rangle$ gives $\langle xy^3, x^3y, xy \rangle = \langle xy \rangle$. See Lemma 5.3 below.

It is known that there exist Stanley decompositions that are not associated in this way with prime filtrations; thus there are also block decompositions that are not associated with subprime decompositions. The standard example, introduced in [11], is $S = K[x, y, z]$, $I = \langle xyz \rangle$,

$$(5.7) \quad S/I \cong K \oplus K[x, y]x \oplus K[y, z]y \oplus K[x, z]z.$$

In this example, the Stanley blocks cannot be ordered in such a way that when added successively to I , they produce monomial ideals. Here

is our version of Soleyman-Jahan's theorem characterizing Stanley decompositions that come from prime filtrations.

5.2. Theorem. *An ordered block decomposition $B_1 \sqcup \cdots \sqcup B_s$ for S/I is subprime if and only if $I \oplus \text{Span} B_1 \oplus \cdots \oplus \text{Span}(B_i)$ is a monomial ideal for each $i = 1, \dots, s$.*

Proof. The \Rightarrow direction has already been proved. For the \Leftarrow direction, the hypothesis implies that $\text{Mon}(I) \sqcup B_1 \sqcup \cdots \sqcup B_k$ is a vector space basis for an ideal I_k for each k , and that these ideals form a filtration. It is equally clear that each $\text{Mon}(I_k) \setminus \text{Mon}(I_{k-1}) = B_k$, so the filtration is subprime. \square

Note that when written as a block decomposition, the blocks of (5.7) are unions of blocks from the elementary decomposition of S/I , just as is the gnomon decomposition. This shows that although the elementary decomposition itself is subprime (Theorem 5.4 below), the result of compressing an elementary decomposition is not in general subprime. But the gnomon decomposition, which is a compression of the elementary decomposition, is subprime (Theorem 5.5).

If B is any block (2.2), and $b_i < \infty$, we define the i th **outer adjacent face** of B to be the block

$$F^i(B) = \begin{bmatrix} b_1 & \cdots & b_{i-1} & b_i + 1 & b_{i+1} & \cdots & b_n \\ a_1 & \cdots & a_{i-1} & b_i + 1 & a_{i+1} & \cdots & a_n \end{bmatrix}.$$

(This is not a subset of B , but is displaced by one step in the direction i from the ordinary i th **outer face** of B .) If $b_i = \infty$ we set $F^i(B) = \emptyset$.

5.3. Lemma. *Let B be a block with inner corner \mathbf{a} , and let I be a monomial ideal disjoint from B . Let $M = B \sqcup \text{Mon}(I)$. Then M spans a monomial ideal (or equivalently, $M = \text{Mon}(J)$ for some monomial ideal J) if and only if the outer adjacent faces of B are contained in I . In this case $B \sqcup \text{Mon}(I) = \text{Mon}(\langle \mathbf{x}^{\mathbf{a}} \rangle) \cup \text{Mon}(I)$ and $J = I + \langle \mathbf{x}^{\mathbf{a}} \rangle$.*

Proof. A set M of monomials spans an ideal if and only if every monomial divisible by (\succeq) an element of M belongs to M . Let $M = B \sqcup \text{Mon}(I)$ and suppose that M spans a monomial ideal. Let $\mathbf{c} \in F^i(B)$ for some i . Then $\mathbf{a} \prec \mathbf{c}$, so $\mathbf{c} \in M$. But $\mathbf{c} \notin B$, so $\mathbf{c} \in I$. Therefore $F^i(B) \subset I$ for each i . Conversely, suppose that $F^i(B) \subset I$ for each i . Since every element of $\langle \mathbf{a} \rangle$ that is not in B is divisible by some element of $F^i(B)$, and therefore belongs to I , we conclude that $\langle \mathbf{a} \rangle \subset M$. Therefore $M = \text{Mon}(I) \cup \text{Mon}(\langle \mathbf{a} \rangle) = \text{Mon}(I + \langle \mathbf{x}^{\mathbf{a}} \rangle)$, so M spans an ideal. \square

We are now ready to prove that elementary and gnomon decompositions (with suitable orderings) are subprime.

5.4. Theorem. *Let $I \subset S = K[x_1, \dots, x_n]$ be a monomial ideal and let $B^1 \sqcup \dots \sqcup B^s$ be the elementary block decomposition of S/I , ordered according to Remark 3.2. Let $B_k = B^{s-k+1}$ for $k = 1, \dots, s$; that is, reverse the ordering of the blocks. Then $B_1 \sqcup \dots \sqcup B_s$ is a subprime decomposition.*

Proof. We work in the superscript notation, restating the criterion of Theorem 5.2 as

$$(5.8) \quad B^k \sqcup \dots \sqcup B^s \sqcup \text{Mon}(I) \quad \text{spans an ideal for each } k = 1, \dots, s.$$

The proof begins with $k = s$ and works backward. Let $A = \{\mathbf{a}^1, \dots, \mathbf{a}^s\}$ be the refined list of inner corners of elementary blocks for S/I , created in step 4 of Algorithm 3.1, in lexicographic order according to Remark 3.2. Since $\mathbf{m} \prec \mathbf{m}'$ implies $\mathbf{m} <_{lex} \mathbf{m}'$, \mathbf{a}^s is a maximal element of A under \prec .

We claim that each outer adjacent face $F^i(B^s)$ is contained in I . Let $\mathbf{c} \in F^i(B^s)$. Then $\mathbf{a}^s \prec \mathbf{c}$ and $\mathbf{c} \notin B^s$. But by Remark 3.3, if $\mathbf{c} \notin I$ then \mathbf{c} belongs to the block of S/I whose inner corner is maximal among inner corners less than \mathbf{c} . But this inner corner must be \mathbf{a}^s , so $\mathbf{c} \in B^s$, a contradiction. Therefore $\mathbf{c} \in I$, proving the claim.

By Lemma 5.3 we conclude that

$$B^s \sqcup I = I + \langle \mathbf{a}^s \rangle = I'$$

is an ideal. Setting $A' = \{\mathbf{a}^1, \dots, \mathbf{a}^{s-1}\}$, we see that A' is the list of inner corners of elementary blocks for S/I' , which has $B^1 \sqcup \dots \sqcup B^{s-1}$ for its elementary decomposition. Now the argument can be repeated, with \mathbf{a}^{s-1} as the maximal element. \square

Lexicographic ordering is not the only ordering that gives a subprime decomposition. It is only necessary to select a maximal element of A , remove it, and repeat.

5.5. Theorem. *The π -gnomon decomposition $B^1 \sqcup \dots \sqcup B^s$ of S/I is subprime, for any π , under the ordering $B_k = B^{s-k+1}$.*

The proof of this theorem actually establishes more, namely, that if Algorithm 4.2 is used to “add a generator” to a subprime block decomposition, the result will be subprime.

Proof. Since any π -gnomon decomposition becomes a standard gnomon decomposition by renumbering the variables, it suffices to prove the standard case. The argument is by induction on the number r of minimal generators of I . If $r = 1$, then

$$\text{Mon}(I) = \begin{bmatrix} \infty & \cdots & \infty \\ m_1 & \cdots & m_n \end{bmatrix}$$

and the gnomon decomposition is

$$\begin{bmatrix} \infty & \cdots & \infty \\ m_1 & \cdots & m_n \\ 0 & \cdots & 0 \end{bmatrix}.$$

In this case Lemma 4.9 shows that when the terms of the gnomon decomposition are added in reverse order to the ideal, an ideal is obtained at each stage.

Now make the induction hypothesis that the theorem is true for all monomial ideals I having r minimal generators. Let I' be a monomial ideal with $r + 1$ generators, and write $I' = I + \langle \mathbf{m} \rangle$, where I has r generators. Suppose that $B^1 \sqcup \cdots \sqcup B^s$ is the gnomon decomposition for S/I . The gnomon decomposition for S/I' is

$$(B^1 \setminus \langle \mathbf{m} \rangle) \sqcup \cdots \sqcup (B^s \setminus \langle \mathbf{m} \rangle) = \bigsqcup_{k=1}^s \left(\bigsqcup_{i=1}^n B^{ki} \right)$$

in the notation of Remark 4.3. We need to prove that for any $p \in \{1, \dots, s\}$ and $q \in \{1, \dots, n\}$,

$$(5.9) \quad (B^{pq} \sqcup \cdots \sqcup B^{pn}) \sqcup \bigsqcup_{k=p+1}^s \left(\bigsqcup_{i=1}^n B^{ki} \right) \sqcup \text{Mon}(I')$$

spans an ideal. From here on we assume p and q are fixed. By the induction hypothesis,

$$(5.10) \quad B^k \sqcup \cdots \sqcup B^s \sqcup \text{Mon}(I)$$

spans an ideal for each $k = 1, \dots, s$. This will be used for $k = p + 1$ and $k = p$. For $k = p + 1$, let J be the ideal spanned by

$$(5.11) \quad B^{p+1} \sqcup \cdots \sqcup B^s \sqcup \text{Mon}(I).$$

For $k = p$, (5.10) together with Lemma 5.3 implies that the outer adjacent faces of B^p belong to J .

Let

$$B^k = \begin{bmatrix} \mathbf{b}^k \\ \mathbf{a}^k \end{bmatrix} \quad \text{and} \quad c_i^k = \max\{a_i^k, m_i\}.$$

By Lemma 4.8 we have

$$\begin{bmatrix} \mathbf{b}^k \\ \mathbf{c}^k \end{bmatrix} = B^k \cap \langle \mathbf{m} \rangle \subset I',$$

so we can include these terms in (5.9) without changing the result, provided we change the last disjoint union to an ordinary union because

at that point there is now an overlap. Thus (5.9) equals

$$(5.12) \quad \left(B^{pq} \sqcup \dots \sqcup B^{pn} \sqcup \begin{bmatrix} \mathbf{b}^p \\ \mathbf{c}^p \end{bmatrix} \right) \\ \sqcup \bigsqcup_{k=p+1}^s \left(\left(\bigsqcup_{i=1}^n B^{ki} \right) \sqcup \begin{bmatrix} \mathbf{b}^k \\ \mathbf{c}^k \end{bmatrix} \right) \cup \text{Mon}(I').$$

By Lemmas 4.1 and 4.8 we have

$$\left(\bigsqcup_{i=1}^n B^{ki} \right) \sqcup \begin{bmatrix} \mathbf{b}^k \\ \mathbf{c}^k \end{bmatrix} = B^k.$$

Since (5.11) spans J and $\text{Mon}(J) \cup \text{Mon}(I') = \text{Mon}(J + I')$,

$$\left(\bigsqcup_{k=p+1}^s B^k \right) \cup \text{Mon}(I')$$

spans $J' = J + I'$. These calculations reduce (5.12) to

$$(5.13) \quad \left(B^{pq} \sqcup \dots \sqcup B^{pn} \sqcup \begin{bmatrix} \mathbf{b}^p \\ \mathbf{c}^p \end{bmatrix} \right) \cup \text{Mon}(J').$$

The induction is completed by considering two cases.

1. If $\mathbf{m}^p \not\preceq \mathbf{b}^p$, then (4.2) applies at the point when the generator \mathbf{m}^p is added during the construction of the gnomon decomposition. So by Remark 4.3, $B^{p1} = B^p$ and $B^{pi} = \emptyset$ for $i = 2, \dots, n$. So (5.13) is trivial unless $p = 1$. But in that case, (5.13) reduces to $B^p \cup \text{Mon}(J')$. The outer adjacent faces of B^p belong to J , and therefore also to J' . So (5.13) spans an ideal.
2. If $\mathbf{m}^p \preceq \mathbf{b}^p$, then Lemma 4.9 implies

$$(5.14) \quad B^{pq} \sqcup \dots \sqcup B^{pn} \sqcup \begin{bmatrix} \mathbf{b}^p \\ \mathbf{c}^p \end{bmatrix} = \begin{bmatrix} b_1^p & \dots & b_{q-1}^p & b_q^p & \dots & b_n^p \\ c_1^p & \dots & c_{q-1}^p & a_q^p & \dots & a_n^p \end{bmatrix}.$$

The outer adjacent faces of this block are contained in those of

$$B^p = \begin{bmatrix} b_1^p & \dots & b_n^p \\ a_1^p & \dots & a_n^p \end{bmatrix},$$

but as noted above, these belong to J' , so again, (5.13) spans an ideal. □

5.6. Remark. To find the subprime filtration associated with a gnomon decomposition, as in Remark 5.1, one should delete any generators that are divisible by each new inner corner that is added. The proof of Theorem 5.5 shows that when a sequence of blocks B^{pn}, \dots, B^{p1} is added to

$\text{Mon}(J')$, their inner corners decrease in \prec . Therefore each new corner added will replace the previous one (and may also remove others). For instance, in Example 4.7 the inner corners $(3, 2, 0)$, $(0, 2, 0)$, $(3, 1, 0)$, $(3, 0, 0)$, $(0, 0, 0)$ are added to the starting ideal with generators $(0, 0, 5)$, $(0, 2, 3)$, $(3, 1, 1)$. The sequences $(3, 2, 0)$, $(0, 2, 0)$ and $(3, 1, 0)$, $(3, 0, 0)$, $(0, 0, 0)$ each arise from a single gnomon, so the inner corners decrease in \prec . The subprime filtration is

$$\begin{aligned} \langle (0, 0, 5), (0, 2, 3), (3, 1, 1) \rangle &\subset \langle (0, 0, 5), (0, 2, 3), (3, 1, 1), (3, 2, 0) \rangle \\ &\subset \langle (0, 0, 5), (3, 1, 1), (0, 2, 0) \rangle \subset \langle (0, 0, 5), (0, 2, 0), (3, 1, 0) \rangle \\ &\subset \langle (0, 0, 5), (0, 2, 0), (3, 0, 0) \rangle \subset \langle (0, 0, 0) \rangle. \end{aligned}$$

When $(0, 2, 0)$ is added, it eliminates both $(3, 2, 0)$, as expected, and also $(0, 2, 3)$, which is not predicted by this remark.

We conclude the section with some technical remarks concerning the usual definition of prime filtration. Recall that a monomial ideal $P \subset S$ is prime if and only if it is generated by a subset Y of the variables x_1, \dots, x_n . Since S is a \mathbb{Z}^n -graded algebra, and monomial ideals in S inherit this multi-grading, quotients of monomial ideals are \mathbb{Z}^n -graded modules. If P is a monomial ideal and $\mathbf{a} \in \mathbb{Z}_{\geq 0}^n$, the symbol $S/P(-\mathbf{a})$ denotes S/P with its \mathbb{Z}^n -grading shifted so that $(S/P)(-\mathbf{a})_{\mathbf{b}} = (S/P)_{\mathbf{b}+\mathbf{a}}$. If $P = \langle Y \rangle$ is a prime ideal, S/P is isomorphic as a vector space to $K[X]$, where X is the complementary set of variables to Y . Then $S/P(-\mathbf{a})$ is isomorphic as a vector space to the Stanley space $K[X]\mathbf{x}^{\mathbf{a}}$. Now (5.1) is defined to be a prime filtration if there exist monomial prime ideals P_i and integer vectors $\mathbf{a}_i \in \mathbb{Z}_{\geq 0}^n$ such that I_i/I_{i+1} is isomorphic as a \mathbb{Z}^n -graded S -module to $S/P_i(-\mathbf{a}_i)$. It is therefore also isomorphic as a vector space to $K[X]\mathbf{x}^{\mathbf{a}}$, but this is not part of the definition.

Using these definitions, Soleyman-Jahan proved Theorem 5.2 above for the case of prime filtrations, with the blocks being Stanley blocks. The difficulty of his proof seems to result from the fact that the association of I_i/I_{i-1} with a particular $S/P_i(-\mathbf{a}_i)$ is only specified by an isomorphism, and not by a concrete construction as we have done. It seems to be the case that if $S/P(-\mathbf{a})$ is isomorphic as a \mathbb{Z}^n -graded S -module to $S/P'(-\mathbf{a}')$, then $P = P'$ and $\mathbf{a} = \mathbf{a}'$. In other words, P_i and \mathbf{a}_i are uniquely determined by I_i and I_{i-1} . This observation (if correct) would bring his treatment closer to ours. However, this uniqueness is only possible if the isomorphism $I_i/I_{i-1} \cong S/P_i(-\mathbf{a}_i)$ is interpreted as an S -module homomorphism; neither algebra homomorphism nor vector space homomorphism is strong enough to achieve the desired

uniqueness. With our approach there is no need to bring in the \mathbb{Z}^n grading or the S -module structure.

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